



However, in a wireless communication system a transmitted signal may still propagate along different paths due to reflection or scattering. Therefore, the signal is received as a superposition of different signal components each propagated along one of the possible paths. As reflected signals will be differently delayed according to their paths compared to each other and compared to the direct signal, each of the signal components of one specific path is interfered by a plurality of other time- and phase-shifted signal components of other paths. If there is no direct sight between the transmitter and the receiver the connection can still be provided by the reflected or scattered signal components, but their interference generally causes disadvantageous effects to the connection performance. Phase-shifted signal components reduce the signal-to-noise ratio of the direct or main signal component and cause power dissipation. Furthermore, the superposition of phase-shifted signals may cause destructive interference between different signal components. These disadvantageous effects reduce the reliability and stability of the decorrelation in the receiver, e.g. increase the bit error rate. Especially when one of the communication partners, e.g. a mobile terminal or phone, is moving these multipath phase shift and interference may change quickly causing a fast fading effect.

To reduce interference of phase shifted signal components several algorithms have been proposed in the past for tracking of the multipath delays. All of them suffer from either high sensitivity to multipath fading and thus reduced performance in such environments, or implementation complexity which is prohibitive for the use in mobile terminals, where size, production cost and power consumption are critical issues. One example for

code-tracking is the so-called early-late gate timing error detector (EL-TED), being a natural solution to a mathematical optimisation problem for nondispersive, additive white gaussian noise channels (AWGN). For a detailed description of prior art see Heinrich Meyr, Marc Moeneclaey and Stefan Fechtel, Digital Communication Receivers: Synchronization, Channel Estimation and Signal Processing, John Wiley and Sons, New York, 1998, Heinrich Meyr and Gerd Ascheid, Synchronization in Digital Communications. Vol. 1, John Wiley & Sons, New York, 1990 and A.J. Viterbi, CDMA - Principles of Spread Spectrum Communication, Addison-Wesley Publishing Company, Reading (MA), 1995.

The EL-TED is able to track echo path delays, as long as the relative delay between two paths is larger than some threshold. If the relative delay is smaller than the threshold, the EL-TED can no longer distinguish the two paths. Furthermore, tracking of weaker paths (lower power) can be disturbed by stronger paths with a relative delay larger than the aforementioned threshold. For an example of an algorithm with significantly higher implementation complexity than the EL-TED, see R.A. Iltis, An EKF-Based Joint Estimator for Interference, Multipath, and Code Delay in a DS Spread-Spectrum Receiver, IEEE Transactions on Communications, Vol. 42, No. 2/3/4, February-April 1994.

In indoor mobile communication scenarios, the path delay of the last identifiable echo path is small compared to outdoor scenarios. All of the signal energy is dispersed around the path with the shortest delay and a sufficient echo path resolution is critical for adequate receiver performance, expressed for instance as the bit-error rate (BER). None of the existing EL-TED algorithms with reasonable complexity is able to track

adjacent path delays which are separated by less than the mentioned threshold. With such a tracking system, the required signal-to-noise ratio to achieve a certain BER is reduced significantly.

5

A system using multipath interference cancellation in the code-tracking loop was presented by Wern-Ho Sheen and Chien-Hsiang Tai, A Noncoherent Tracking Loop With Diversity and Multipath Interference Cancellation for Direct-Sequence Spread-Spectrum Systems, IEEE Transactions on Communications, Vol. 46, No. 11, November 1998. There, a group of rake fingers is tracked jointly. In each finger, respreading and recoding is performed after the normal rake correlation in order to yield signals which model the interference for each multipath. These interference terms can then be subtracted from the incoming signal in order to improve the tracking performance, inherently by allowing a finger assignment where adjacent fingers are placed around one chip duration apart. This scheme suffers from high implementation complexity, expressed by the respreading and recoding necessary in each finger.

Therefore, it is an object of the present invention to provide a method and/or a rake receiver for code-tracking in CDMA communication systems of low complexity and cost.

Another object is to provide a method and/or rake receiver with stable and reliable code-tracking as well as improved performance, reducing the disadvantages of the prior art.

Surprisingly simple, these objects of the invention are already achieved by the method according to claim 1 and the rake receiver according to claim 16.

According to the present invention an electromagnetic signal of a CDMA communication system is received. The received signal is generated from the transmitted signal by a superposition of signal components of different signal paths. The transmitted signal is preferably spread and coded by a pseudo-noise bit sequence. In particular, different signal paths are created by scattering, diffraction and/or reflection of the signal. Therefore, the signal components may be time- and phase-shifted to each other and having different amplitudes.

The received signal is preferably filtered and digitised by an analog-to-digital (A/D) converter. The digitised signal is distributed or conveyed to a plurality ( $N$ ) of receiver fingers which are constituents of a so-called rake receiver.

Furthermore, the digitised signal is distributed or conveyed to a detection and a synchronising streams or branches, preferably in each of the receiver fingers. The signal is preferably distributed by a digital interpolator/decimator.

In the detection and synchronisation streams the signal is preferably correlated with the complex-conjugate of the pseudo-noise bit sequence used in the transmitter. The decorrelated signal in the synchronisation stream is multiplied with reconstructed transmitted symbols.

The method further comprises a step of reducing interference of at least one signal component of signal path  $j$  with the signal component of signal path  $i$  by an interference reduction device of a receiver finger which

is assigned to signal path  $i$ , where  $i \neq j$ .

Preferably reducing of interference comprises a subtraction of an interference signal from the decorrelated digitised signal. The interference signal is preferably calculated using the complex weight coefficients of at least one or all signal paths  $j$  ( $j \neq i$ ). Preferably the subtraction takes place on symbol rate  $1/T$ . Preferably, interference reduction is provided in each rake finger. Advantageously, interference reduction provides a stable and reliable code-tracking of low complexity and cost. Furthermore, the method and device according to the invention yield an improved estimation of the timing delay  $\tau$ .

In a preferred embodiment of the invention the decorrelation by the pseudo-noise sequence takes place in a timing error detector which includes a first and second correlators. The digitised signal is transmitted to the first correlator and the time-shifted signal is transmitted to the second correlator. The signals are correlated, preferably multiplied with the complex-conjugate of the pseudo-noise sequence in the first and second correlator, respectively, and the two resulting output signals are subtracted. The decorrelated signal is preferably multiplied with reconstructed transmitted symbols.

In a further embodiment the real part of the complex output signal of the interference reduction device is filtered. The output signal of the filter provides an estimated time delay  $\hat{\tau}^{(i)}$  for the signal component of the respective signal path  $i$ .

The invention is described in detail hereinafter by means of preferred embodiments and reference is made to the attached drawings.

5 **Brief description of the Figures**

It is shown in

Fig. 1 a block diagram of a rake receiver with  $N$  receiver fingers.

Fig. 2 a block diagram of a receiver finger  $i$  of a rake receiver with a conventional coherent early-late timing error detector,

Fig. 3 a block diagram of a receiver finger  $i$  of the rake receiver of Fig. 1 with an early-late timing error detector according to a first embodiment of the invention,

Fig. 4 a diagram of the detector characteristic of a conventional coherent early-late timing error detector,

Fig. 5 a diagram of the detector characteristic of a conventional coherent early-late timing error detector for two signal components  $i$  and  $j$ ,

Fig. 6 a block diagram of a receiver finger  $i$  of a rake receiver according to a second embodiment of the invention and

Fig. 7 a block diagram of a receiver finger  $i$  of a rake receiver according to a third embodiment of the invention.

**Detailed Description of the Invention**

The block diagrams of Fig. 1, 3, 5 and 6 can also be read as flow charts of the method according to the invention. Same or equivalent items of the conventional timing error detector 102' shown in Fig. 2 and of the timing error detector 102 according to the present

invention are denoted by the same numerals.

Fig. 1 presents the general structure of code-tracking loops for a plurality of receiver fingers 1, 2, ..., N of a rake receiver 17 for a direct sequence CDMA communication system. For direct sequence CDMA the signal is spread by a factor of  $2^L$  in the transmitter and multiplied by a code sequence before modulation, yielding chips of a duration  $T_c$ . For the preferred embodiment quadrature phase shift keying (QPSK) and a pseudo-noise (PN) code sequence 112 are used. Also binary phase shift keying (BPSK) or any other linear modulation scheme can be applied.

A digital spread spectrum rake receiver 17 including a digital interpolator/decimator 16 is shown in Fig. 1. A wirelessly transmitted signal 10 is received by the device 100 and is fed through a pulse matched filter 12. The filter 12 is a root raised-cosine pulse filter, matched to the transmission filter (not shown) which is also of root raised-cosine shape in the transmitter (not shown). As best seen in Fig. 2 and 3 an analog-to-digital (A/D) converter 14 generates samples  $z(t = n \cdot T_s)$  15 of the pulsed matched-filtered signal 13 at equidistant times  $t$  which are integer multiples  $n$  of the sample time  $T_s$  and feeds them to the digital interpolator/decimator 16. The digital interpolator/decimator 16 generates intermediate samples  $z(t = k \cdot T_c + \hat{t})$  110 of the same signal at equidistant times  $t$  which are integer multiples  $k$  of the chip duration  $T_c$  shifted by estimated timing instants  $\hat{t}$  for each of the  $N$  identical rake fingers. The signal is sampled at a rate which is higher than the chip rate, typical oversampling factors are in the range of four to eight. The digital interpolator/decimator 16 interpolates



between the samples, as sampling rate  $1/T_s$  and chip rate  $1/T_c$  can be different, and performs a sampling rate conversion from the sampling rate to the chip rate.

5        The interpolated signal 110 is fed to the detection stream or branch, where a correlation of the signal 110 with the complex-conjugate pseudo-noise (PN) code 112 in each of the  $N$  fingers takes place. In the detection stream path delay and phase shift of each signal component is equalised and the signal amplitude is optimised. For reconstruction of the transmitted symbols phase shift elimination and amplitude optimisation is achieved by multiplication of the signal with the complex-conjugate path weight  $c_k^{(i)*}$ . This reconstruction yields an enhancement of the strong signal paths and an additional weakening of already weak signal paths. After this reconstruction the signals of all fingers are added up in a combiner (not shown).

20        Furthermore, the signal is fed from the digital interpolator decimator 16 to  $N$  code-tracking loops, where exemplary the code-tracking loops of finger 1, finger 2 and finger  $N$  are denoted 101, 201 and 301, respectively, as seen in Fig. 1. Each of the  $N$  code-tracking loops includes a timing error detector (TED) which generates error signals that are substantially proportional to the remaining path delay error in the signal. The timing error detector of the first code-tracking loop or first finger is denoted 102. The error signals are lowpass-filtered in a loop filter to yield estimates for the respective path delay, again for each finger, which are then fed back to a respective input of the interpolator/decimator 16, closing the tracking loop.

35        A conventional coherent early-late timing error

detector 102' embedded in a code-tracking loop for one of  
 the fingers  $i$  of a rake receiver is depicted in Fig. 2.  
 The timing error detector and other components downstream  
 of the digital interpolator/decimator 16 is replicated  
 5 for each rake finger  $i \in \{1, 2, \dots, N\}$ , whereas the  
 interpolator/decimator 16 is used jointly. The samples  
 $z(t = n \cdot T_s)$  15 of the incoming signal are interpolated and  
 decimated at the estimated timing delay  $\hat{\tau}^{(i)}$  for the  
 respective signal path assigned to the finger  $i$ . A first  
 10 and second data streams are generated, the first stream  
 is fed to the detection stream or path, more specifically  
 to a correlator 120. It is noted that the detection  
 stream is not shown in full detail. The second data  
 stream is fed to the synchronisation stream or path, more  
 15 specifically to a correlator 121 and a shifter component ( $z$   
 '1) 123. In the first stream the samples  $z(t = k \cdot T_c + \hat{\tau})$   
 110 are correlated with the pseudo-noise code 112 in the  
 correlator 120. In the second or synchronisation stream  
 shifted samples  $z(t = k \cdot T_c - T_c/2 + \hat{\tau})$  111, spaced by -  
 20  $T_c/2$  with respect to the samples in the detection stream,  
 are correlated with the pseudo-noise code in the  
 correlator 121. Shifted samples 111 are further shifted  
 by  $+T_c$  in the component 123 to feed samples  $z(t = k \cdot T_c +$   
 $T_c/2 + \hat{\tau})$  in the correlator 122. The two correlators 121,  
 25 122 of the synchronisation stream yielding early and late  
 estimates 114, 113 of the transmitted symbols. Estimates  
 113, 114 are substantially raised-cosine functions. If  
 synchronisation of the samples is perfect, i.e. estimated  
 time delay  $\hat{\tau}$  is close to the real delay of the respective  
 30 signal path, early and late estimates are spaced by  $\pm$   
 $T_c/2$  from the maximum of said raised-cosine function,  
 respectively, and the difference of early and late  
 estimates provided in step 124 is close to zero. When

estimated delay  $\hat{\tau}$  moves away from the real delay, difference of early and late estimates become finite yielding a delay compensation in the digital interpolator/estimator. Thus the tracking loop provides a  
5 synchronisation or locking loop for the digital signal. According to a next step 125 the difference signal 117 is multiplied with reconstructed transmitted symbols 115. Reconstructed transmitted symbols 115 are generated by complex-conjugately multiplying pilot symbols, e.g. of  
10 training sequences or symbol decisions with estimates of the channel phase or the channel phasor, but other methods of symbol reconstruction can also be applied. The real part of the resulting samples is extracted by circuit 126, now on symbol rate  $1/T$ , as error signal  $\tilde{x}$   
15 118. Error signal 118 is fed through the loop filter 103 with lowpass characteristic to yield the estimate for the timing delay  $\hat{\tau}^{(i)}$ . The general method of timing error detection is known in the art and e.g. described in Heinrich Meyr, Marc Moeneclaey and Stefan Fechtel,  
20 Digital Communication Receivers: Synchronization, Channel Estimation and Signal Processing, John Wiley and Sons, New York, 1998, A.J. Viterbi, CDMA - Principles of Spread Spectrum Communication, Addison-Wesley Publishing Company, Reading (MA), 1995, A. Wilde, Extended tracking  
25 range delay-locked loop, Proceedings of the International Conference on Communications, Seattle, USA, 1995 and Volker Aue and Gerhard P. Fettweis, A Non-Coherent Tracking Scheme for the RAKE Receiver That Can Cope With Unresolvable Multipath, Proceedings of the International  
30 Conference on Communications, Vancouver, Canada, 1999.

Herewith, said documents are incorporated as subject matter of the present disclosure by reference.

Fig. 3 shows one of the identical receiver fingers  
35 with an early-late timing error detector 102 with

adaptive interference reduction at symbol rate according to the invention, embedded in a code-tracking loop for one of the fingers  $i$  in the rake receiver 17. Ideally interference is nearly cancelled as result of signal subtraction 130 in an interference reduction device 131. The chiprate processing part is equivalent to that of the conventional timing error detector 102' shown in Fig. 2.

According to the present invention the error signal coming out of the timing error detector 102 and going into the loop filter 103, is modified in the inventive timing error detector adaptively by computing the interference from adjacent multipaths. After multiplication of the reconstructed symbols in the current finger  $i$  according to step 125, interference from other fingers  $j$  ( $j \neq i$ ) is subtracted in step 130 from an intermediate signal 116 on symbol rate  $1/T$ . For computing the interference channel state information such as path amplitudes and phases as well as path locations (relative to the current path  $i$ ) are used. The path locations are inherently available from the tracking loops of the other rake fingers  $j$  ( $j \neq i$ ), while the path amplitudes and the phases can be derived from the channel tracking units (not shown) of the other fingers.

As an alternative to interpolating and decimating the data signal 15 for each rake finger, the pseudo-noise code 112 can also be shifted in order to guarantee the correct alignment between data samples and pseudo-noise code samples. Furthermore, due to normally very slowly varying path delays, delay tracking of all fingers can be implemented in time-sharing using only one dedicated tracking structure in an alternative embodiment of the invention.

A typical detector characteristic for the conventional coherent timing error detector, also called S-curve, is depicted in Fig. 4. The S-curve depends on the pulse filter used in the transmitter and the corresponding matched filter 12 in the receiver. Fig. 4 shows the expected value of the error signal  $x^{(i)}$  for an additive white gaussian noise channel and a raised-cosine pulse-matched filter autocorrelation function. The current path and as such the current rake finger is denoted by index  $i$  or exponent  $(i)$ . In a multipath environment, the timing error detection is influenced by adjacent multipaths lying everywhere but in the zero-crossings of the S-curve. As an example, Fig. 5 further shows an S-curve of a second path  $j$ , causing a weaker signal component than path  $i$  and being time-shifted approximately by one chip duration  $T_c$  with respect to path  $i$ . It is clearly seen that at the usual location of operation of the tracking loop for finger  $i$  which is the origin in the diagram, an additive interference from finger  $j$  will degrade the performance of the timing error detection. The interference depends on the distance of the two paths  $i$  and  $j$ , on the instantaneous amplitude and also on the phase relation.

If the current channel conditions are known, either perfect or by means of estimates, i.e. path amplitudes and phases as well as path delays are known, the interference from other fingers are estimated and subtracted from the timing error signal of the current finger. The rake receiver has a total of  $N$  fingers, and complex instantaneous channel tap phasors, which include amplitude and phase information, are denoted by the complex vector  $c_k = [c_k^{(1)} \dots c_k^{(N)}]$ . Furthermore, the path delays

are given by  $\tau = [\tau_k^{(1)} \dots \tau_k^{(N)}]$ , where index  $k$  defines the current symbol for the path delays and the phasors.

The advantages of the inventive timing error detection will be clear from Fig. 4 and 5 representing the expected value of the error signal  $E[x]$  for a non interfered path  $i$  (Fig. 4) and exemplary for another scenario where two paths  $i$  and  $j$  are interfering each other (Fig. 5). As shown in Fig. 4 the S-Curve is a function of the path delay  $\tau$ , therefore, shall be denoted by  $S(\tau)$  and is determined by the structure of the early-late timing error detector. On the x- and y-axis the timing delay  $\tau$  in units of  $T_c$  and the expected value  $E[x]$  are plotted, respectively. As long as the path  $i$  is not interfered the conventional timing error detector is unbiased, i.e. the value of the S-curve at the origin is zero, as shown in Fig. 4. Fig. 5 depicts the signal components for two different signal paths  $i$  and  $j$ , where the amplitude of component  $j$  is lower than the one of component  $i$ . The timing error detector for finger  $i$  is now biased, meaning that the superimposed contribution from path  $j$  results in a nonzero value of the S-curve for finger  $i$  at the origin. In advantageous contrast to that the timing error detector for finger  $i$  according to the invention subtracts interference of path  $j$  from the signal component of path  $i$  such that reliable and stable path delay estimation and tracking in the tracking loop is provided for each signal component.

To provide the above described interference reduction the S-curve of the transmission system is stored in an interference computation module 132 being part of the interference reduction device 131. The

interference computation module 132 receives the path delays 133  $\hat{\tau}_k^{(i)}$  and  $\hat{\tau}_k^{(j)}$  for the assigned signal path  $i$  and for all other paths  $j \neq i$ , respectively. Further, module 132 receives the complex path weights  $c_k^{(j)}$  134 of all other paths  $j \neq i$  than the assigned signal path  $i$ .

In particular, for path  $i$ , the interference estimates of all other fingers are subtracted from the intermediate signal  $y_k^{(i)}$  116 according to

$$\tilde{x}_k^{(i)} = \text{Re}\{y_k^{(i)} - \sum_{j=1, j \neq i}^M c_k^{(i)*} c_k^{(j)} S(\hat{\tau}_k^{(i)} - \hat{\tau}_k^{(j)})\} \quad (1)$$

The first factor  $c_k^{(i)*}$  accounts for the complex-conjugate of the maximum ratio coefficient used in finger  $i$ . The second factor  $c_k^{(j)}$  stands for the channel tap phasor which the interfering path is subject to on the channel itself. The product of both is the total weight the interfering path is subject to in finger  $i$ , and multiplied with the S-curve at the estimated correct location,  $\hat{\tau}_k^{(i)} - \hat{\tau}_k^{(j)}$ , it results in the interference contribution of finger  $j$  to finger  $i$ . The storage of  $S$  in the module 132 may be either by means of a mathematical function (e.g. an approximation by  $\sin(\tau)/\tau$ ) such that  $S(\hat{\tau}_k^{(i)} - \hat{\tau}_k^{(j)})$  is calculated for each argument or  $S$  is stored by means of a look-up table such that the values of the function  $S$  are interpolated from the table values.

In the exemplary embodiment depicted in Fig. 1 and 3 the inventive timing error detection is applied for all receiver fingers  $i \in \{1, 2, \dots, N\}$ . Thus the invention

yields for each finger  $i$  an interference cancelled or at least reduced error signal  $\tilde{x}_k^{(i)}$  118 and an interference reduced estimated timing delay  $\hat{\tau}_k^{(i)}$  after the loop filter 103. Therefore, a significantly improved tracking is provided for all paths compared to prior art.

In another preferred embodiment, shown in Fig. 6, the real part of the intermediate signal  $y_k^{(i)}$  116 is extracted before reducing the interference. For this embodiment interference computation 132' is simplified by computing only the real part of the sum in Eq. (1) and it becomes

$$\tilde{x}_k^{(i)} = x_k^{(i)} - \sum_{j=1, j \neq i}^M \operatorname{Re}\{c_k^{(i)*} c_k^{(j)}\} S(\hat{\tau}_k^{(i)} - \hat{\tau}_k^{(j)}) \quad (2)$$

Figure 7 shows an alternative early-late structure, where the order of correlation and subtraction has been reversed. This is allowable due to the linear nature of the correlation operation and yields a simplified structure with only one correlator 121'.